

A Logical Approach to Discrete Math

Sets

Set enumeration

List each element in the set.

Example: $\{0, 2, 4, 6, 8\}$

Set comprehension

$\{x : t \mid R : E\}$

↑ ↑ ↑ ↑
Body

Range

Optional type

Dummy variable

Example: $\{x : \mathbb{Z} \mid 0 \leq x < 5 : 2 \cdot x\}$

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Set comprehension

$$\{x : t \mid R : E\}$$

Set types

The type of the set is not the type of the body.

If E has type $t1$, the set has type $set(t1)$.

Type $t1$ is not necessarily the same as type t .

Example

$\{x : \mathbb{Z} \mid 0 \leq x < 5 : 2 \cdot x\}$ has type $set(\mathbb{Z})$.

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Set comprehension

$$\{x : t \mid R : E\}$$

From comprehension to enumeration

$$\{i \mid 0 < i < 4 : i\} = \{1, 2, 3\}$$

$$\{i \mid 5 \leq i \leq 8 : 2 \cdot i\} = \{10, 12, 14, 16\}$$

$$\{i \mid 9 \leq i \leq 17 \wedge \text{even}.i : i\} = \{10, 12, 14, 16\}$$

$$\{x, y \mid 2 \leq x \leq 3 \leq y \leq 4 : x^y\} = \{2^3, 2^4, 3^3, 3^4\}$$

$$\{x \mid 0 \leq x < 3 : x \cdot y\} = \{0 \cdot y, 1 \cdot y, 2 \cdot y\}$$

$$\{x \mid 0 \leq x < 0 : x \cdot y\} = \{\} \quad \text{the empty set}$$

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(11.2) **Axiom, Enumeration:**

$$\{e_0, e_1, \dots, e_{n-1}\} = \{x \mid x = e_0 \vee x = e_1 \vee \dots \vee x = e_{n-1} : x\}$$

From enumeration to comprehension

$$\{6, 17\} = \{x \mid x = 6 \vee x = 17 : x\}$$

$$\{a, b, c\} = \{x \mid x = a \vee x = b \vee x = c : x\}$$

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(11.3) **Axiom, Set membership:** Provided $\neg occurs('x', 'F')$,
 $F \in \{x \mid R : E\} \equiv (\exists x \mid R : F = E)$

Examples

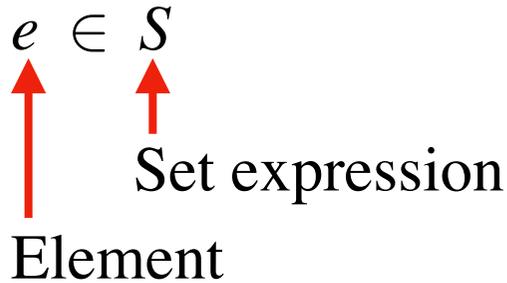
$$2 \in \{i \mid 0 < i < 4 : i\} \equiv (\exists i \mid 0 < i < 4 : 2 = i)$$

$$13 \in \{i \mid 5 \leq i \leq 8 : 2 \cdot i\} \equiv (\exists i \mid 5 \leq i \leq 8 : 13 = 2 \cdot i)$$

A Logical Approach to Discrete Math

(11.3) **Axiom, Set membership:** Provided $\neg occurs('x', 'F')$,
 $F \in \{x \mid R : E\} \equiv (\exists x \mid R : F = E)$

\in as a binary infix operator

$e \in S$

↑ Set expression
↑ Element

Example: $25 \in \{13, 25\} \equiv true$

\in as a function

$elementOf(e, S)$ returns true if e is an element of S .

Example: $elementOf(25, \{13, 25\})$ returns *true*.

In this example, the type of the function is

$elementOf: \mathbb{Z} \times set(\mathbb{Z}) \rightarrow \mathbb{B}$

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(11.3) **Axiom, Set membership:** Provided $\neg occurs('x', 'F')$,
 $F \in \{x \mid R : E\} \equiv (\exists x \mid R : F = E)$

Theorem: $e \in \{e\}$

Proof

$$\begin{aligned} & e \in \{e\} \\ = & \langle (11.2) \text{ Enumeration} \rangle \\ & e \in \{x \mid x = e : x\} \\ = & \langle (11.3) \text{ Set membership} \rangle \\ & (\exists x \mid x = e : e = x) \\ = & \langle (8.14) \text{ One-point rule} \rangle \\ & (e = x)[x := e] \\ = & \langle \text{Textual substitution} \rangle \\ & e = e \quad \text{which is reflexivity of } = \quad // \end{aligned}$$

A Logical Approach to Discrete Math

A Theory of Sets

- (11.2) **Axiom, Enumeration:** $\{e_0, e_1, \dots, e_{n-1}\} = \{x \mid x = e_0 \vee x = e_1 \vee \dots \vee x = e_{n-1} : x\}$
- (11.3) **Axiom, Set membership:** Provided $\neg \text{occurs}('x', 'F')$,
 $F \in \{x \mid R : E\} \equiv (\exists x \mid R : F = E)$
- (11.4) **Axiom, Extensionality:** $S = T \equiv (\forall x \mid : x \in S \equiv x \in T)$
- (11.4.1) **Axiom, Empty set:** $\emptyset = \{x \mid \text{false} : E\}$
- (11.4.2) $e \in \emptyset \equiv \text{false}$
- (11.4.3) **Axiom, Universe:** $\mathbf{U} = \{x \mid : x\}$, $\mathbf{U} : \text{set}(t) = \{x : t \mid : x\}$
- (11.4.4) $e \in \mathbf{U} \equiv \text{true}$, for $e : t$ and $\mathbf{U} : \text{set}(t)$
- (11.5) $S = \{x \mid x \in S : x\}$

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(11.5.1) **Axiom, Abbreviation:** For x a single variable, $\{x \mid R\} = \{x \mid R : x\}$

(11.6) Provided $\neg occurs('y', 'R')$ and $\neg occurs('y', 'E')$,

$$\{x \mid R : E\} = \{y \mid (\exists x \mid R : y = E)\}$$

(11.7) $x \in \{x \mid R\} \equiv R$

R is the characteristic predicate of the set.

(11.7.1) $y \in \{x \mid R\} \equiv R[x := y]$ for any expression y

(11.9) $\{x \mid Q\} = \{x \mid R\} \equiv (\forall x \mid : Q \equiv R)$

(11.10) $\{x \mid Q\} = \{x \mid R\}$ is valid iff $Q \equiv R$ is valid.

(11.11) **Methods for proving set equality $S = T$:**

(a) Use Leibniz directly.

(b) Use axiom Extensionality (11.4) and prove the (9.8) Lemma

$v \in S \equiv v \in T$ for an arbitrary value v .

(c) Prove $Q \equiv R$ and conclude $\{x \mid Q\} = \{x \mid R\}$.

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(11.4) **Axiom, Extensionality:** $S = T \equiv (\forall x | : x \in S \equiv x \in T)$

(9.8) $(\forall x | R : true) \equiv true$

Proving set equality

Use (11.4) and (9.8) to prove that two sets, S and T , are equal.

Lemma: A small theorem used as a step in the main theorem.

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(11.4) **Axiom, Extensionality:** $S = T \equiv (\forall x | : x \in S \equiv x \in T)$

(9.8) $(\forall x | R : true) \equiv true$

Example

Prove (11.5) $S = \{x | x \in S : x\}$ without a lemma.

Proof

$$\begin{aligned} & S = \{x | x \in S : x\} \\ = & \langle (11.4) \text{ Extensionality} \rangle \\ & (\forall v | : v \in S \equiv v \in \{x | x \in S : x\}) \\ = & \dots \\ & (\forall v | : true) \\ = & \langle (9.8) \rangle \\ & true \quad // \end{aligned}$$

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Example

Prove (11.5) $S = \{x \mid x \in S : x\}$ with a lemma.

The (9.8) lemma

Let v be an arbitrary element, and prove that

$$v \in S \equiv v \in \{x \mid x \in S : x\}$$

Proof

$$\begin{aligned} & v \in \{x \mid x \in S : x\} \\ = & \langle (11.3) \text{ Set membership} \rangle \\ & (\exists x \mid x \in S : v = x) \\ = & \langle (9.20.1) \text{ Existential double trading} \rangle \\ & (\exists x \mid v = x : x \in S) \\ = & \langle (8.14) \text{ One-point rule} \rangle \\ & v \in S \quad // \end{aligned}$$

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(11.4) **Axiom, Extensionality:** $S = T \equiv (\forall x | : x \in S \equiv x \in T)$

(9.8) $(\forall x | R : true) \equiv true$

The main proof

$$S = \{x \mid x \in S : x\}$$

$$= \langle (11.4) \text{ Extensionality} \rangle$$

$$(\forall v | : v \in S \equiv v \in \{x \mid x \in S : x\})$$

$$= \langle \text{The (9.8) lemma} \rangle$$

$$(\forall v | : true)$$

$$= \langle (9.8) \rangle$$

$$true \quad //$$

From now on, you may omit the main proof.

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R is the characteristic predicate of the set.

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(11.11) **Methods for proving set equality $S = T$:**

(a) Use Leibniz directly.

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 $v \in S \equiv v \in T$ for an arbitrary value v .

(c) Prove $Q \equiv R$ and conclude $\{x \mid Q\} = \{x \mid R\}$.

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Operations on sets.

(11.12) **Axiom, Size:** $\#S = (\sum x \mid x \in S : 1)$

(11.13) **Axiom, Subset:** $S \subseteq T \equiv (\forall x \mid x \in S : x \in T)$

(11.14) **Axiom, Proper subset:** $S \subset T \equiv S \subseteq T \wedge S \neq T$

(11.15) **Axiom, Superset:** $T \supseteq S \equiv S \subseteq T$

(11.16) **Axiom, Proper superset:** $T \supset S \equiv S \subset T$

(11.17) **Axiom, Complement:** $v \in \sim S \equiv v \in \mathbf{U} \wedge v \notin S$

(11.18) $v \in \sim S \equiv v \notin S$, for v in \mathbf{U}

(11.19) $\sim\sim S = S$

(11.20) **Axiom, Union:** $v \in S \cup T \equiv v \in S \vee v \in T$

(11.21) **Axiom, Intersection:** $v \in S \cap T \equiv v \in S \wedge v \in T$

(11.22) **Axiom, Difference:** $v \in S - T \equiv v \in S \wedge v \notin T$

(11.23) **Axiom, Power set:** $v \in \mathcal{P}S \equiv v \subseteq S$

A Logical Approach to Discrete Math

$$(11.12) \quad \text{Axiom, Size: } \#S = (\sum x \mid x \in S : 1)$$

Example

$$U : \{a, b, c, d, e, f\}$$

$$S : \{b, e, f\}$$

$$\#S = 1 + 1 + 1 = 3$$

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(11.13) **Axiom, Subset:** $S \subseteq T \equiv (\forall x \mid x \in S : x \in T)$

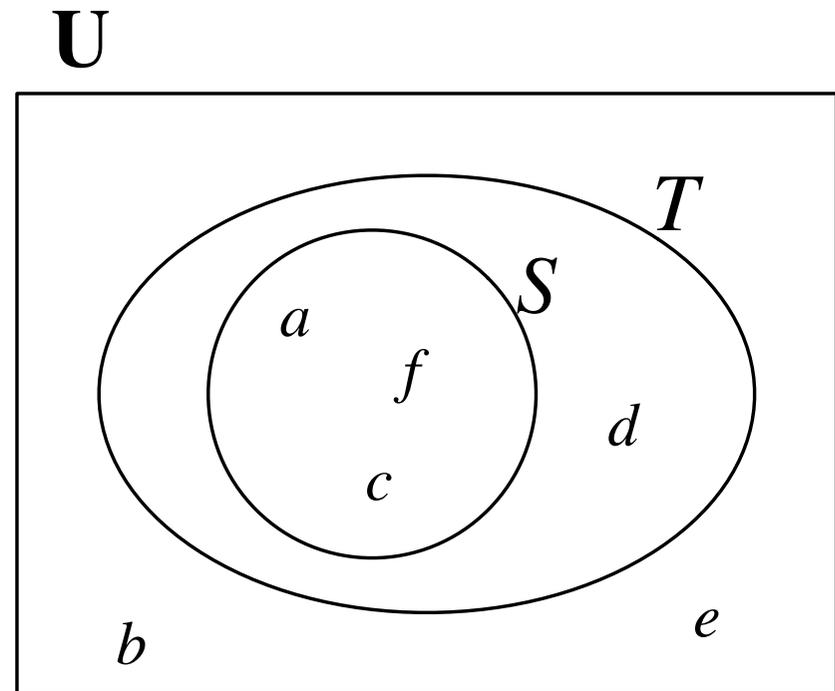
Example

$U : \{a, b, c, d, e, f\}$

$S : \{a, c, f\}$

$T : \{d, f, c, a\}$

$S \subseteq T$



A Logical Approach to Discrete Math

(11.13) **Axiom, Subset:** $S \subseteq T \equiv (\forall x \mid x \in S : x \in T)$

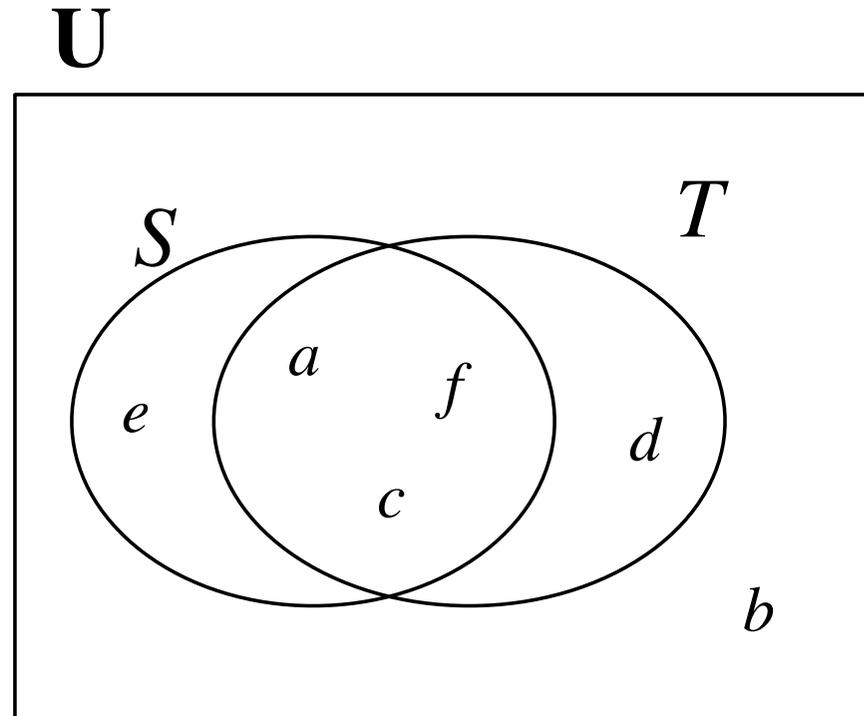
Example

$U : \{a, b, c, d, e, f\}$

$S : \{a, c, f, e\}$

$T : \{d, f, c, a\}$

$S \not\subseteq T$



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(11.13) **Axiom, Subset:** $S \subseteq T \equiv (\forall x \mid x \in S : x \in T)$

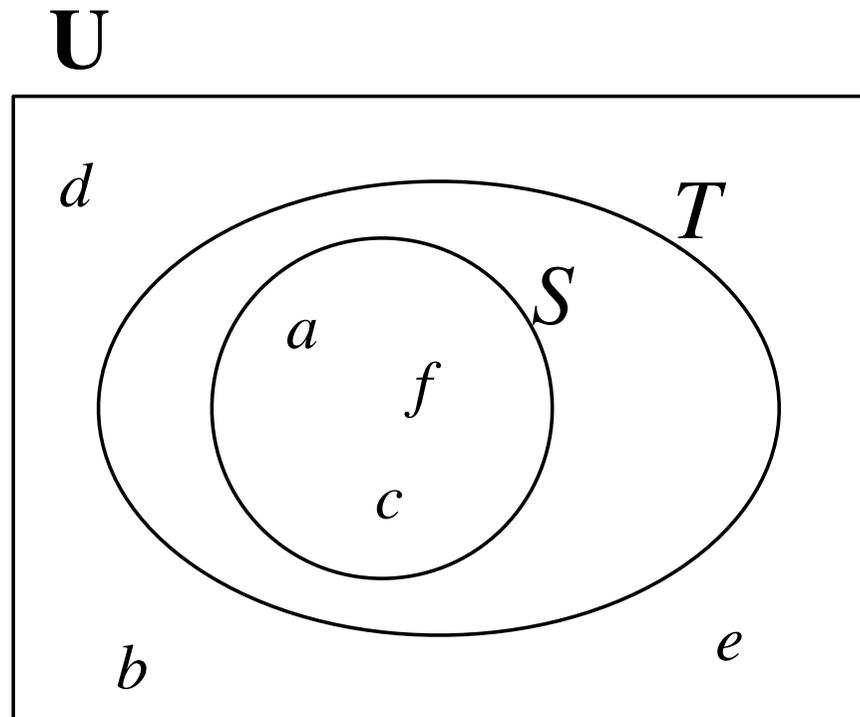
Example

$U : \{a, b, c, d, e, f\}$

$S : \{a, c, f\}$

$T : \{f, c, a\}$

$S \subseteq T$



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(11.14) **Axiom, Proper subset:** $S \subset T \equiv S \subseteq T \wedge S \neq T$

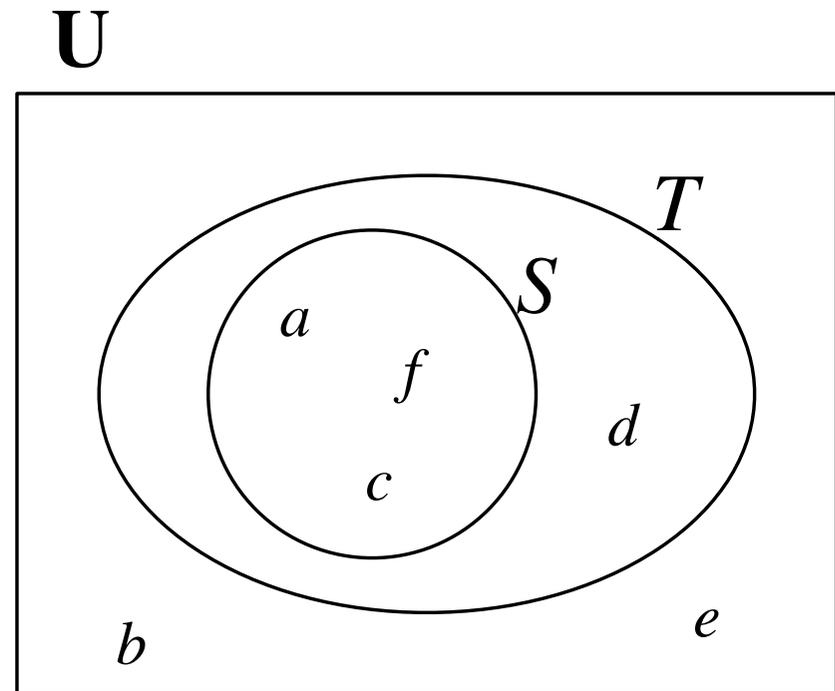
Example

$U : \{a, b, c, d, e, f\}$

$S : \{a, c, f\}$

$T : \{d, f, c, a\}$

$S \subset T$



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(11.14) **Axiom, Proper subset:** $S \subset T \equiv S \subseteq T \wedge S \neq T$

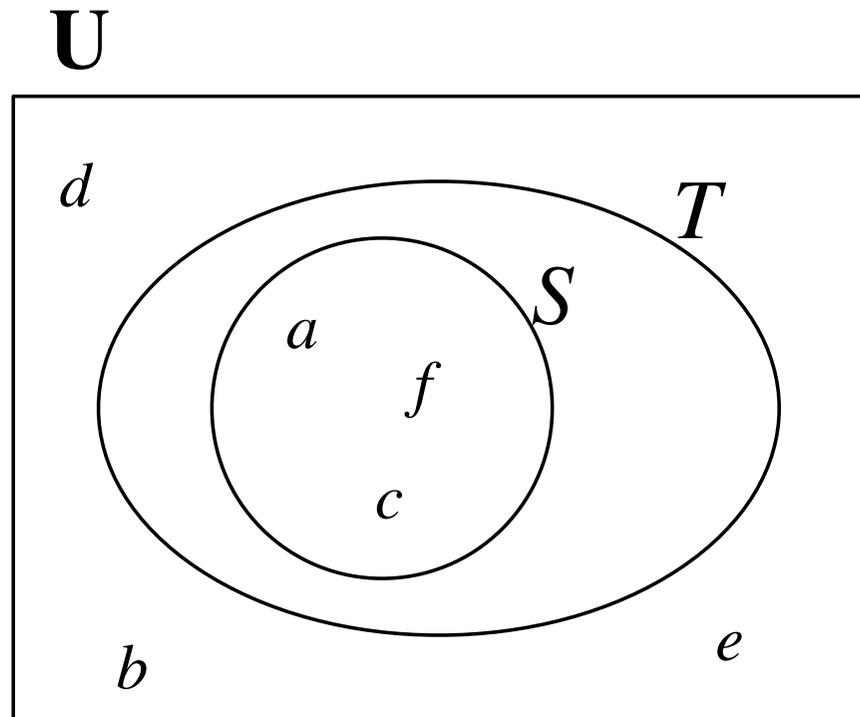
Example

$U: \{a, b, c, d, e, f\}$

$S: \{a, c, f\}$

$T: \{f, c, a\}$

$S \not\subseteq T$



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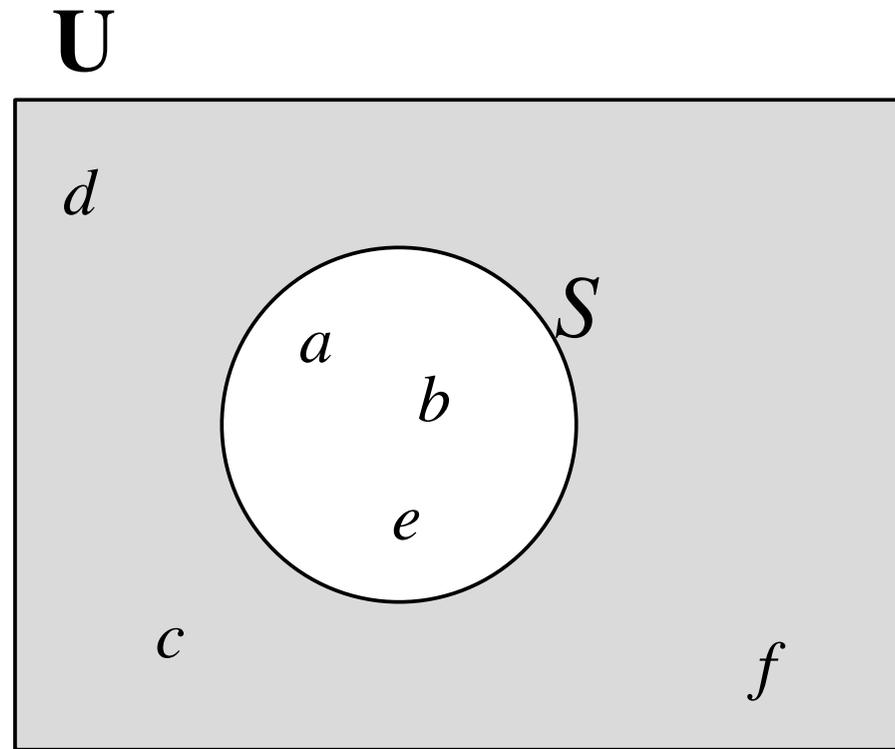
(11.17) **Axiom, Complement:** $v \in \sim S \equiv v \in \mathbf{U} \wedge v \notin S$

Example

$\mathbf{U} : \{a, b, c, d, e, f\}$

$S : \{a, b, e\}$

$\sim S = \{c, d, f\}$



A Logical Approach to Discrete Math

$$(11.20) \quad \text{Axiom, Union:} \quad v \in S \cup T \equiv v \in S \vee v \in T$$

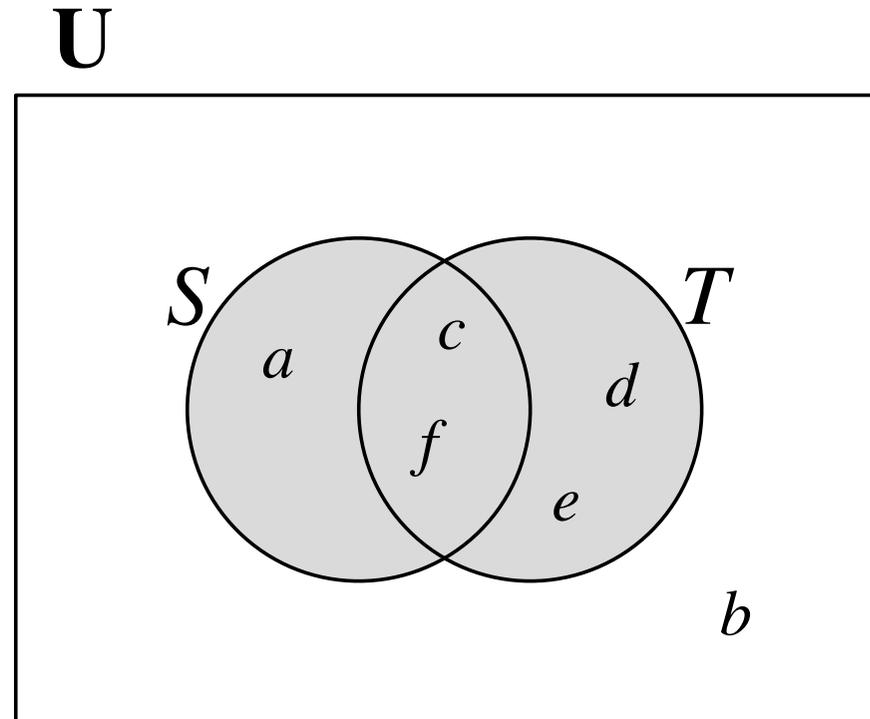
Example

$$U: \{a, b, c, d, e, f\}$$

$$S: \{a, c, f\}$$

$$T: \{c, d, e, f\}$$

$$S \cup T = \{a, c, d, e, f\}$$



A Logical Approach to Discrete Math

(11.21) **Axiom, Intersection:** $v \in S \cap T \equiv v \in S \wedge v \in T$

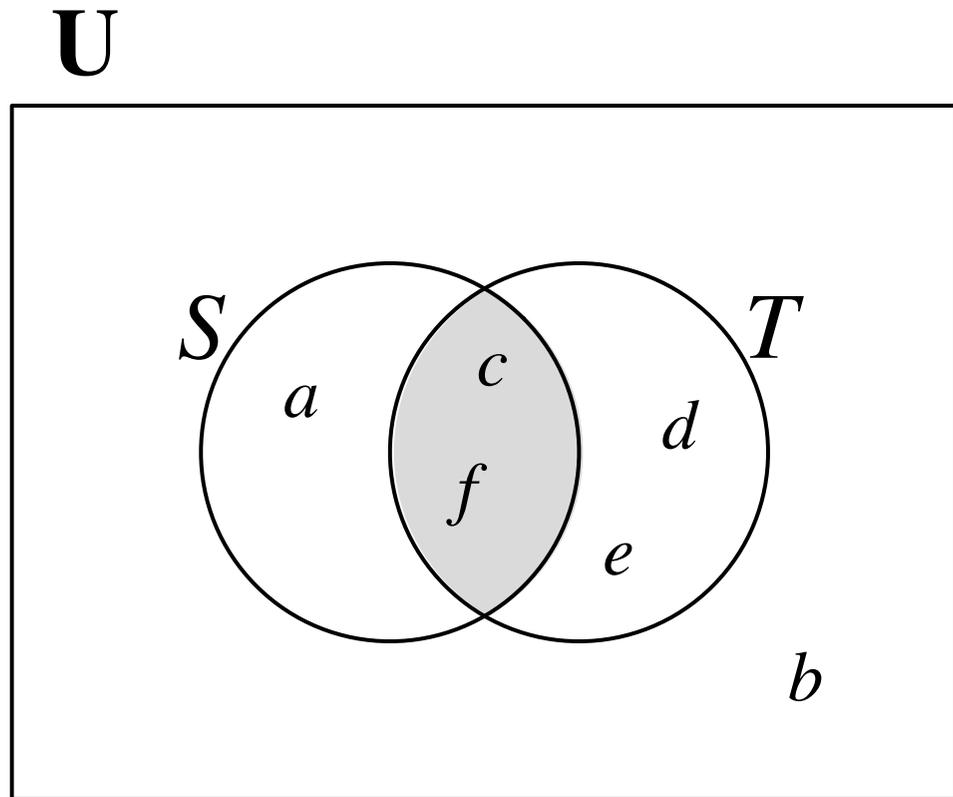
Example

$U: \{a, b, c, d, e, f\}$

$S: \{a, c, f\}$

$T: \{c, d, e, f\}$

$S \cap T = \{c, f\}$



A Logical Approach to Discrete Math

(11.22) **Axiom, Difference:** $v \in S - T \equiv v \in S \wedge v \notin T$

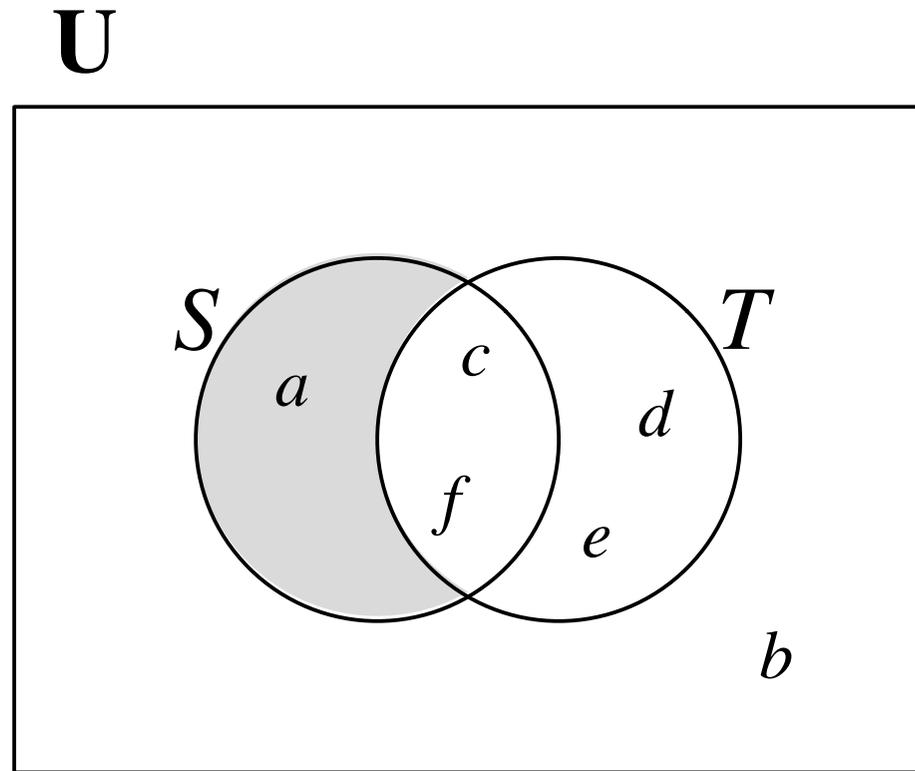
Example

$U: \{a, b, c, d, e, f\}$

$S: \{a, c, f\}$

$T: \{c, d, e, f\}$

$S - T = \{a\}$



A Logical Approach to Discrete Math

(11.23) **Axiom, Power set:** $v \in \mathcal{P}S \equiv v \subseteq S$

Example

$S : \{a, b\}$

$\mathcal{P}S = \{\{\}, \{a\}, \{b\}, \{a, b\}\}$

Example

$S : \{a, b, c\}$

$\mathcal{P}S = \{\{\}, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$

A Logical Approach to Discrete Math

TABLE 11.1. TYPES OF SET EXPRESSIONS IN THEORY $set(t)$

Expression	Example (with types)	Type of result
Empty set, universe, variable	\emptyset or U or S	$set(t)$
Set enumeration	$\{e_1:t, \dots, e_n:t\}$	$set(t)$
Set comprehension	$\{x \mid R:\mathbb{B} : E:t\}$	$set(t)$
	$\{x:t \mid R:\mathbb{B}\}$	$set(t)$
Set membership	$x:t \in S:set(t)$	\mathbb{B}
Set equality	$S:set(t) = T:set(t)$	\mathbb{B}
Set size	$\# S:set(t)$	\mathbb{N}
$\subset, \supset, \subseteq, \supseteq$	$S:set(t) \subseteq T:set(t)$	\mathbb{B}
Complement	$\sim S:set(t)$	$set(t)$
$\cup, \cap, -$	$S:set(t) \cup T:set(t)$	$set(t)$
Power set	$(\mathcal{P} S):set(t)$	$set(set(t))$

A Logical Approach to Discrete Math

- (11.24) **Definition.** Let E_s be a set expression constructed from set variables, \emptyset , \mathbf{U} , \sim , \cup , and \cap . Then E_p is the expression constructed from E_s by replacing:
 \emptyset with *false*, \mathbf{U} with *true*, \cup with \vee , \cap with \wedge , \sim with \neg .
The construction is reversible: E_s can be constructed from E_p .

Example

$E_s : S \cup \sim S$ This S is a set expression.

$E_p : S \vee \neg S$ This S is a boolean expression.

Example

$E_s : \emptyset$ This is a set expression.

$E_p : \textit{false}$ This is a boolean expression.

A Logical Approach to Discrete Math

(11.25) **Metatheorem.** For any set expressions E_s and F_s :

(a) $E_s = F_s$ is valid iff $E_p \equiv F_p$ is valid,

(b) $E_s \subseteq F_s$ is valid iff $E_p \Rightarrow F_p$ is valid,

(c) $E_s = \mathbf{U}$ is valid iff E_p is valid.

Example

$S \cap \sim S = \emptyset$ is valid iff $S \wedge \neg S \equiv \textit{false}$ is valid.

Example

$S \cap T \subseteq S$ is valid iff $S \wedge T \Rightarrow S$ is valid.

A Logical Approach to Discrete Math

Prove (11.51) $S - \emptyset = S$

Proof

$$\begin{aligned} S - \emptyset &= S \\ &= \langle(11.49)\rangle \\ S \cap \sim \emptyset &= S \end{aligned}$$

By Metatheorem (11.25a) $S \cap \sim \emptyset = S$ is valid
iff $S \wedge \neg false \equiv S$ is valid

$$\begin{aligned} S \wedge \neg false & \\ &= \langle(3.13)\rangle \\ S \wedge true & \\ &= \langle(3.39)\rangle \\ S & // \end{aligned}$$

A Logical Approach to Discrete Math

CAUTION: You cannot use a metatheorem in a proof hint.

The following step is illegal.

Prove (11.51) $S - \emptyset = S$

Proof

$$S - \emptyset = S$$

$$= \langle (11.49) \rangle$$

$$S \cap \sim \emptyset = S$$

$$= \langle (11.25a) \rangle \leftarrow \text{Illegal}$$

$$S \wedge \neg \text{false} \equiv S$$

A Logical Approach to Discrete Math

Quantifying union and intersection

(11.74.1) **Definition:** $v \in (\cup x \mid R : E) \equiv (\exists x \mid R : v \in E)$

(11.75.1) **Definition:** $v \in (\cap x \mid R : E) \equiv (\forall x \mid R : v \in E)$

Quantifying union \cup

\cup is symmetric: (11.26) $S \cup T = T \cup S$

\cup is associative: (11.27) $(S \cup T) \cup U = S \cup (T \cup U)$

\cup has an identity: (11.30) $S \cup \emptyset = S$

Therefore, \cup is an abelian monoid and can be quantified.

A Logical Approach to Discrete Math

(11.76) **Axiom, Partition:** Set S partitions T if

(i) the sets in S are pairwise disjoint and

(ii) the union of the sets in S is T , that is, if

$$(\forall u, v \mid u \in S \wedge v \in S \wedge u \neq v : u \cap v = \emptyset) \wedge (\cup u \mid u \in S : u) = T$$

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Example

$$T : \{a, b, c, d, e, f\}$$

$$S : \{\{a, c\}, \{b, e, f\}, \{d\}\}$$

S partitions T .

A Logical Approach to Discrete Math

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Example

$$T : \{a, b, c, d, e, f\}$$

$$S : \{\{a, c\}, \{b, e, f\}, \{d, e\}\}$$

S does not partition T because $\{b, e, f\} \cap \{d, e\} \neq \emptyset$.

A Logical Approach to Discrete Math

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Example

$$T : \{a, b, c, d, e, f\}$$

$$S : \{\{a, c\}, \{e, f\}, \{d\}\}$$

S does not partition T because $\{a, c\} \cup \{e, f\} \cup \{d\} \neq T$.

A Logical Approach to Discrete Math

Bags

Same as sets, but duplicates allowed.

$$\begin{aligned}\{x:\mathbb{N} \mid -2 \leq x \leq 2 : x^2\} &= \{4, 1, 0, 1, 4\} \\ \{x:\mathbb{N} \mid -2 \leq x \leq 2 : x^2\} &= \{4, 1, 0\}\end{aligned}$$

A Logical Approach to Discrete Math

Bags

(11.79) **Axiom, Membership:** $v \in \{x \mid R : E\} \equiv (\exists x \mid R : v = E)$

(11.80) **Axiom, Size:** $\#\{x \mid R : E\} = (\Sigma x \mid R : 1)$

(11.81) **Axiom, Number of occurrences:**

$$v\#\{x \mid R : E\} = (\Sigma x \mid R \wedge v = E : 1)$$

(11.82) **Axiom, Bag equality:** $B = C \equiv (\forall v \mid : v\# B = v\# C)$

(11.83) **Axiom, Subbag:** $B \subseteq C = (\forall v \mid : v\# B \leq v\# C)$

(11.84) **Axiom, Proper subbag:** $B \subset C = B \subseteq C \wedge B \neq C$

(11.85) **Axiom, Union:** $B \cup C = \{v, i \mid 0 \leq i < v\# B + v\# C : v\}$

(11.86) **Axiom, Intersection:**

$$B \cap C = \{v, i \mid 0 \leq i < v\# B \downarrow v\# C : v\}$$

(11.87) **Axiom, Difference:**

$$B - C = \{v, i \mid 0 \leq i < v\# B - v\# C : v\}$$